

Exact Static Compaction of Sequential Circuit Tests Using Branch-and-Bound and Search State Registration

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The paper presents a new method for static compaction of sequential circuit tests that are divided into independent test sequences. We propose an exact method based on the branch-and-bound approach. The search space for the algorithm is efficiently pruned at each step by determining the set of essential vectors, removing faults and sequences implementing the domination relationships and identifying equivalent search states. The method is tested on a large number of benchmark test sets. Experiments show that, unlike previous approaches, this method is capable of finding and proving globally optimal results for all the compaction benchmarks.

Static compaction of test sets consisting of independent test sequences

Basic definitions and problem formulation

A test set T consists of test sequences $t_i \in T$, $i = 1, \dots, n$. Each sequence t_i contains in turn L_i test vectors. We refer to L_i as the *test length* of sequence t_i . The set of faults f_j , $j = 1, \dots, m$ detected by T is denoted by F . Total test length of test set T can be viewed as a sum

$$\sum_{i=1}^n L_i.$$

Test set T consisting of n faults and m test sequences can be viewed as the following matrix, where $t_{si,fj}$ is equal to k if sequence s_i covers fault f_j at the k -th vector and zero if sequence s_i does not cover fault f_j .

$$T = \begin{pmatrix} t_{f_1,s_1} & t_{f_2,s_1} & \dots & t_{f_n,s_1} \\ t_{f_1,s_2} & t_{f_2,s_2} & \dots & t_{f_n,s_2} \\ \dots & \dots & \dots & \dots \\ t_{f_1,s_m} & t_{f_2,s_m} & \dots & t_{f_n,s_m} \end{pmatrix}$$

If we select k vectors from sequence s_i then all the faults $\{f_j : k \geq t_{si,fj} > 0\}$ are said to be covered by these vectors. Our task is to cover all the faults by selecting the minimal number of vectors. As shown in [1], this task is an NP-complete problem.

Example

Consider the test set that consists of three test sequences s_1 , s_2 and s_3 . Sequence s_1 consists of 4 test vectors covering fault f_2 at the 3-rd vector and f_1 at the 4-th vector. Sequence s_2 consists of three test vectors covering f_1 at the first vector and f_3 at the third vector. Finally, sequence s_3 consists of four test vectors covering f_2 at the first vector, f_3 at the second vector and f_4 at the fourth vector.

It can be seen that the minimal solution would be selecting sequence s_3 and the first vector of sequence s_2 .

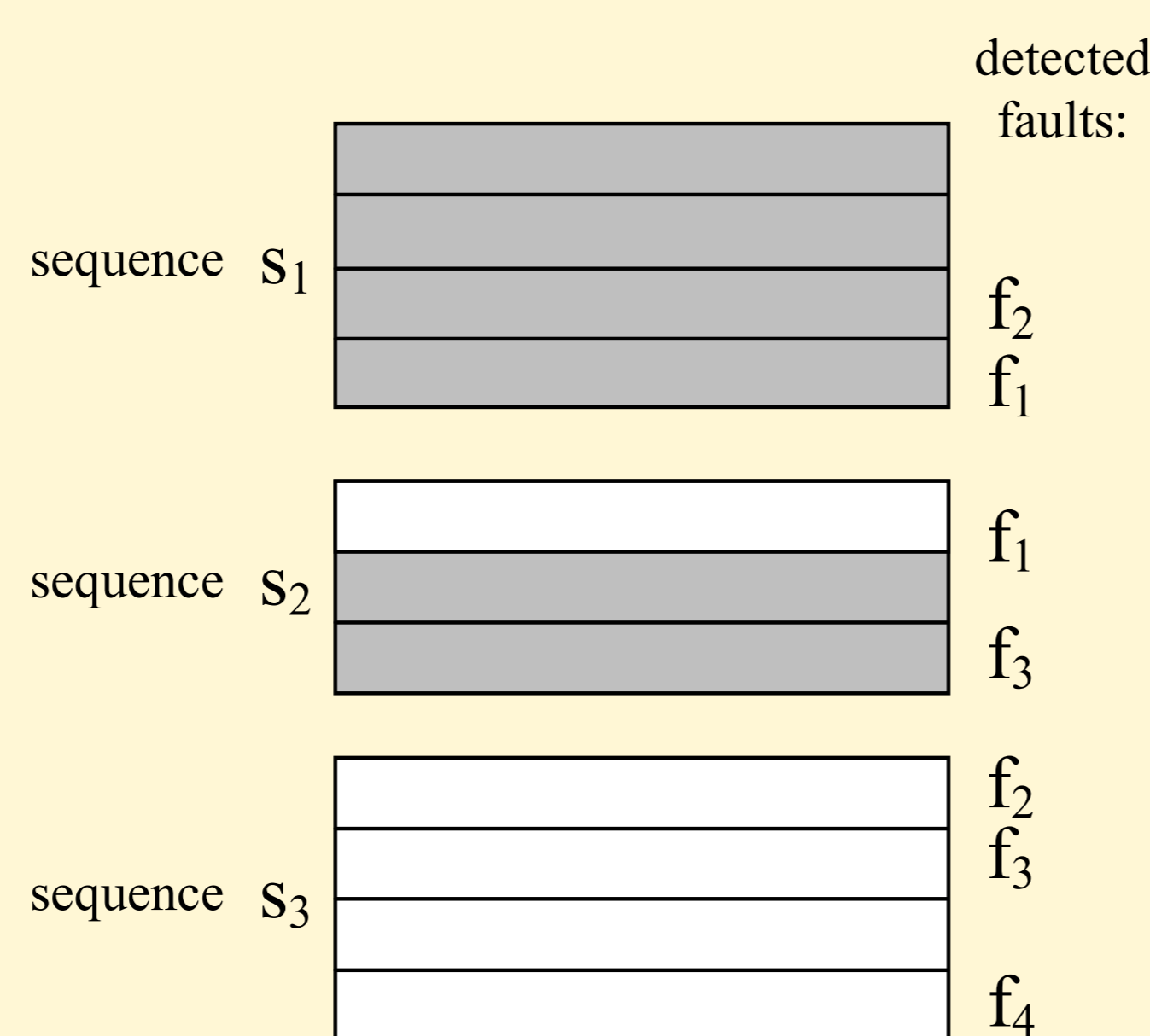


Figure 1. Test set example

Main steps

- *Essential vectors* are detected and removed from the test sequences. If fault f_j is detected by the k -th vector of test sequence s_i and is not detected by any other sequence then k first vectors of sequence s_i are called essential.

- During *removing dominated faults*, column f_a will be removed from matrix T if there exists another column f_b , where

$$\forall_{i=1}^m t_{s_i,f_b} \neq 0 \Rightarrow t_{s_i,f_a} \neq 0, \quad t_{s_i,f_b} \geq t_{s_i,f_a}.$$

- Another type of implications is *removing dominating sequences*. A row corresponding to sequence s_b is said to be a dominating sequence of s_a iff

$$\forall_{j=1}^n t_{s_b,f_j} \neq 0 \Rightarrow t_{s_a,f_j} \neq 0, \quad t_{s_a,f_j} \leq t_{s_b,f_j}.$$

- Current technique implements a branch-and-bound algorithm, which uses depth-first approach for the decision tree traversal. The search is improved by discarding decision combinations equivalent to previously traversed ones.

Algorithm:

```

Compaction()
{
  Repeat
  {
    Select essential vectors
    If current bound is exceeded
      Return
    If all faults covered
      Save result, set new bound, and return
    Remove dominating sequences
    Remove dominated faults
  }
  Until exist essential vectors
  Make sequence selection
  If rank of the selection lower than that of previous
    selection in the decision -tree
  { Return }
  If the selection is lower than current bound
  {
    If all faults covered
      Save result, set new bound, and return
    Else
      Call Compaction()
  }
  Return
}

```

Experimental results

circuit	Circuit size # faults	Initial test set # seq. # vec.	Essential test # seq. # vec.	Result in [1] # seq. # vec.	Result in [2] # seq. # vec.	Current approach # seq. # vec. time, s
s344.g	322	19 141	6 49	10 66	10 66	10 66 0,01
s349.g	330	19 144	9 75	11 84	11 84	11 84 0,01
s420.g	453	33 797	5 325	8 333	8 333	8 333 0,01
s510.g	550	37 989	4 146	7 237	7 263	7 237 0,05
s820.g	816	38 669	13 335	14 347	14 347	14 347 0,03
s838.g	929	37 1323	5 323	11 473	11 482	11 473 0,05
s938.g	929	37 1323	5 323	11 473	11 482	11 473 0,05
s953.g	1053	75 1099	26 447	32 539	32 539	32 539 0,15
s967.g	1038	72 1223	27 606	31 669	31 669	31 669 0,14
s1238.g	1327	123 1554	62 956	72 1007	74 1004	74 1004 16,4
s1269.g	1309	52 450	23 198	29 245	29 245	29 245 3,39
s1512.g	1281	52 772	12 261	14 289	15 294	14 289 0,09
s3271.g	3206	132 2529	43 1047	50 1532	53 1210	50 1178* 27,60
s3330.g	2866	108 2028	39 1018	44 1067	45 1070	43 1067 0,51
s3384.g	3360	58 888	17 327	22 410	22 410	22 410 0,49
s4863.g	4666	112 1533	31 666	42 746	42 749	41 745* 2,41
s5378.g	4603	71 919	37 464	41 493	42 493	42 493 0,58
s6669.g	6506	64 592	29 240	36 303	36 301	36 301 1,12
s38417.g	27733	95 1617	22 588	31 697	31 698	30 684* 2,75
s38584.g	36303	271 8065	95 3416	- -	106 3812	105 3806* 341,9
s641.h	465	78 306	31 150	36 170	36 170	36 170 0,05
s838.h	929	52 675	7 297	12 310	12 310	12 310 0,04
s938.h	929	52 675	7 297	12 310	12 310	12 310 0,04
s1196.h	1214	189 509	105 322	109 339	109 337	109 337 0,16
s3271.h	3206	61 984	12 327	22 489	19 489	19 489 0,40
s4863.h	4666	105 376	55 250	57 257	57 256	57 256 0,45
s35932.h	38448	376 1712	6 188	13 244	11 242	11 242 291,7
s1196.s	1214	297 613	195 365	200 376	199 375	199 375 0,27
s1494.s	1490	160 1787	94 1118	100 1140	99 1140	99 1140 0,35

Conclusions

- The use of implications at each iteration to considerably reduce the search space for the compaction algorithm
- Branch-and-Bound algorithm with identification of equivalent search states. This requires decision ordering.
- The first approach to find and prove globally optimal results for all the ISCAS'89 test sets
- Fast compaction: up to 341.9 s (for the s38584.g test set) on a 366 MHz UltraSPARC computer

References

- [1] F. Corno, P. Prinetto, M. Rebaudengo, M. Sonza Reorda, "New static compaction techniques of test sequences for sequential circuits". *Proc. ED&TC*, 1997, pp.37-43.
- [2] J. Raik, A. Jutman, R. Ubar, "Fast Static Compaction of Test Sequences Using Implications and Greedy Search" *Proc. of ETW*, Stockholm, Sweden, May 29 – June 1, 2001, pp. 207-209.



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